Semidefinite optimization for polynomials in noncommuting variables



Networks & Optimization Algorithm & Complexity

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What are we up to?

- Generalize polynomial optimization over scalar variables
- Want to optimize polynomials evaluated in matrices

What do we need?

- Polynomials in noncommuting variables EAT #TEA
- Approximation technique using semidefinite programs A



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What do I need it for?

- Applications in quantum physics ³/₂
 - quantum chemistry: ground state electronic energy of atoms
 - quantum theory: upper bounds for violation of Bell inequalities
 - quantum information: multi prover games/quantum correlation
- Application in systems control
 - Systematic strategy to compute stabilizing feedback for closed loop systems







What is polynomial optimization?

▶ $p \in \mathbb{R}[\underline{x}]$ polynomial

Find

$$p_{min} = \min_{\underline{a} \in \mathbb{R}^n} p(\underline{a})$$

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What is polynomial optimization?





What is polynomial optimization?





Example

A matrix *M* is copositive if $p_{min} \ge 0$ for $p = \sum_{i,j} M_{ij} x_i^2 x_j^2$.

Problem

Calculating p_{min} is in general NP-hard \bigotimes

- Find a way to make it easier \rightarrow approximation
- Involves sums of squares and semidefinite programs 😳

$$\begin{array}{ll} \max \ \langle \boldsymbol{C}, \boldsymbol{X} \rangle \\ \text{s.t.} \ \langle \boldsymbol{A}_{j}, \boldsymbol{X} \rangle = \boldsymbol{b}_{j}, \quad j = 1, \dots, m \\ \boldsymbol{X} \succeq \boldsymbol{0} \end{array}$$

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$$\max \langle C, X \rangle$$

s.t. $\langle A_j, X \rangle = b_j, \quad j = 1, \dots, m$
 $X \succeq 0$

 $\leftarrow \text{linear function}$

Optimization of a linear function

$$\begin{array}{c|c} \max \langle C, X \rangle \\ \text{s.t.} \langle A_j, X \rangle = b_j, \quad j = 1, \dots, m \\ X \succeq 0 \end{array} \quad \leftarrow \text{ affine space}$$

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Optimization of a linear function over an affine space



$$\begin{array}{c|c} \max \langle \mathcal{C}, X \rangle \\ \text{s.t.} \langle \mathcal{A}_j, X \rangle = \mathcal{b}_j, \quad j = 1, \dots, m \\ X \succeq 0 \qquad \qquad \leftarrow \text{psd matrix} \end{array}$$

Optimization of a linear function over an affine space intersected with the set of positive semidefinite matrices:



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spectrahedron

Optimization of a linear function over an affine space intersected with the set of positive semidefinite matrices: a spectrahedron





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Optimization of a linear function over an affine space intersected with the set of positive semidefinite matrices: a spectrahedron



poly in log(1/ ε) for precision ε

 Essentially solvable in polynomial time using interioir point algorithms, e.g. SeDuMi, SDPT3, SDPA, Mosek,...

NC polynomial optimization



Idea: Replace \underline{a} with $a_i \in \mathbb{R}$ by \underline{A} with A_i symmetric matrices

- Model NC polynomials
 - ▶ Polynomials in noncommuting variables $\underline{X} = (X_1, ..., X_n)$
 - ► Like usual polynomials, only difference $X_1X_2 \neq X_2X_1$ EAT \neq TEA

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- Evaluation in symmetric matrices

$$p = 1 + 2X_1^2 + X_2X_1 - X_1X_2,$$

$$p(A) = 1_0 + 2A_1^2 + A_2A_1 - A_1A_2$$

$$\bullet \underline{A} = (A_1, A_2) \in (S\mathbb{R}^{s \times s})^2 \qquad \int \mathcal{P}(\underline{A}) - \mathbf{I}_s + 2A_1 + A_2A_1 - A_1A_2$$

NC polynomial optimization



Idea: Replace \underline{a} with $a_i \in \mathbb{R}$ by \underline{A} with A_i symmetric matrices

- Model NC polynomials
 - ▶ Polynomials in noncommuting variables $\underline{X} = (X_1, ..., X_n)$
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- Evaluation in symmetric matrices

$$p = 1 + 2X_1^2 + X_2X_1 - X_1X_2, \underline{A} = (A_1, A_2) \in (S\mathbb{R}^{s \times s})^2$$

$$p(\underline{A}) = \mathbf{1}_s + 2A_1^2 + A_2A_1 - A_1A_2$$

NC polynomial optimization

 $p_{min} = \min_{(\varphi,\underline{A})} \{ \langle \varphi, p(\underline{A})\varphi \rangle \mid \|\varphi\| = 1 \}$

▶ p_{min} is the smallest eigenvalue p(<u>A</u>) can attain over all <u>A</u> We can add polynomial constraints like $g(\underline{A}) \succeq 0$ to define a region where we want to optimize p

Where is the SDP?



We can reformulate our nc optimization problem

$$\boldsymbol{p}_{min} = \min_{(\varphi,\underline{A})} \{ \langle \varphi, \boldsymbol{p}(\underline{A})\varphi \rangle \mid \|\varphi\| = 1 \}$$

using nc sums of squares

This will turn out to be a semidefinite program

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But first, let's look at applications of nc polynomial optimization

Application: Quantum Chemistry



Compute ground state energy of atoms

- Molecule of N electrons that can occupy M orbitals
- Each orbital associated with creation/anihilation operators a_i^{\dagger}, a_i
- Pairwise interaction described by h_{ijkl}

$$\begin{split} \min_{(a,a^{\dagger},\varphi)} &\left\langle \varphi, \sum_{ijkl} h_{ijkl} a_i^{\dagger} a_j^{\dagger} a_k a_l \varphi \right\rangle \\ \text{s.t.} \quad \|\varphi\| &= 1 \\ &\left\{ a_i, a_j \right\} = \left\{ a_i^{\dagger}, a_j^{\dagger} \right\} = 0 \\ &\left\{ a_i^{\dagger}, a_j \right\} = \delta_{ij} \\ &\left(\sum_i a_i^{\dagger} a_i - N \right) \varphi = 0 \end{split}$$







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$$\begin{array}{c} \min_{(a,a^{\dagger},\varphi)} \left\langle \varphi, \sum_{ijkl} h_{ijkl} a_{i}^{\dagger} a_{j}^{\dagger} a_{k} a_{l} \varphi \right\rangle \\ \text{s.t.} \quad \|\varphi\| = 1 \\ \quad \{a_{i}, a_{j}\} = \{a_{i}^{\dagger}, a_{j}^{\dagger}\} = 0 \\ \quad \{a_{i}^{\dagger}, a_{j}\} = \delta_{ij} \\ \quad (\sum_{i} a_{i}^{\dagger} a_{i} - N)\varphi = 0 \end{array} \qquad \leftarrow \begin{array}{c} \langle \varphi, p(a, a^{\dagger}) \varphi \rangle \\ \leftarrow \|\varphi\| = 1 \\ \text{additional constraints} \end{array}$$







Application: Systems Control

 \blacktriangleright Linear closed loop system with unknown feedback ${\cal G}$

 $\begin{array}{c} u \\ \hline \\ v \\ \hline \\ v \\ \hline \\ feedback \mathcal{G} \\ \hline \end{array}$

Math. System $\vec{x}(t) = \mathcal{A}\vec{x}(t) + \mathcal{B}\vec{u},$ $\vec{y}(t) = \mathcal{C}\vec{x}(t)$

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► Goal Find *G* which stabilizes the system

Application: Systems Control





Math. System $\vec{x}(t) = \mathcal{A}\vec{x}(t) + \mathcal{B}\vec{u},$ $\vec{y}(t) = \mathcal{C}\vec{x}(t)$

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Lyapunov's idea can be extended to our problem: Riccati equations

Application: Systems Control







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- Lyapunov's idea can be extended to our problem: Riccati equations
- Optimization problem is first a feasibility problem
- Can be refined by optimizing a specific singular value
- For a uniform strategy to get \mathcal{G} we have to work free of dimensions

Application: Quantum Correlations

- Two separated systems $A = M_1 \cup \cdots \cup M_n$ and $B = M_{n+1} \cup \cdots \cup M_N$
- Measurements of *M_i* described by operators *E_i* performed on a joint quantum state φ
- ► Correlations between *A* and *B*: Joint probabilities $P(i,j) = \langle \varphi, E_i E_j \varphi \rangle$



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Application: Quantum Correlations

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- Measurements of M_i described by operators E_i performed on a joint quantum state φ
- Correlations between *A* and *B*: Joint probabilities $P(i, j) = \langle \varphi, E_i E_j \varphi \rangle$
- Violation of Bell inqualities
 - Linear combination of (joint) probabilities
 - Get inequalities by considering classical random variables
 - Want to find violations using quantum setup

$$\begin{array}{l} \max_{(E,\varphi)} \left\langle \varphi, \sum_{i,j} c_{ij} E_i E_j \varphi \right\rangle & \leftarrow \left\langle \varphi, p(\underline{E}) \varphi \right\rangle \\ \text{s.t.} \quad \|\varphi\| = 1 & \leftarrow \|\varphi\| = 1 \\ E_i E_j = \delta_{ij} \text{ for } i, j \in M_k \\ \sum_{i \in M_k} E_i = 1 & \\ [E_i, E_j] = 0 \text{ for } i \in A, j \in B \end{array} \right\} \begin{array}{l} \text{Measurement} \\ A / B \text{ separated} \end{array}$$



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