Communication using quantum entanglement: benefits and limitations



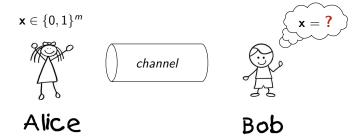


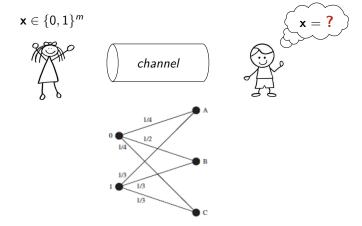
Teresa Piovesan

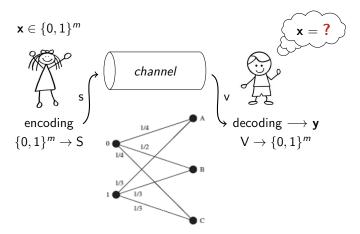
CWI Scientific Meeting Amsterdam, 13 May 2016

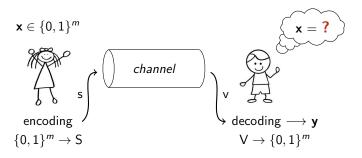






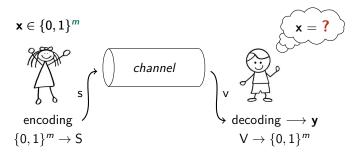






GOAL:

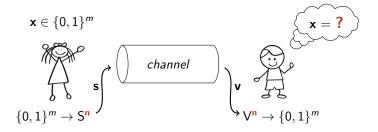
- $\mathbf{y} = \mathbf{x}$ with high probability



GOAL:

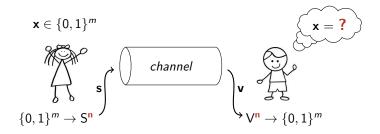
- $\mathbf{y} = \mathbf{x}$ with high probability
- maximize m

Block coding



• Encoding x into a sequence of channel inputs can be more efficient

Block coding

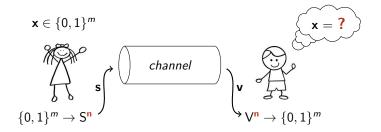


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Shannon (1948):

Number of bits that can be theoretically transmitted per channel use

Block coding

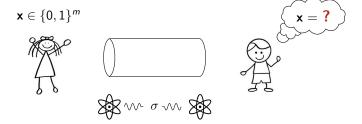


• Encoding x into a sequence of channel inputs can be more efficient

Shannon (1948):

Number of bits that can be theoretically transmitted per channel use

• error probability goes to zero as $\mathbf{n} \to \infty$

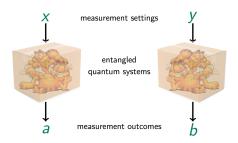


Quantum entanglement

Entanglement is one of the most striking features of quantum mechanics Measurements on *entangled* quantum systems can give outcomes that are non-classically correlated

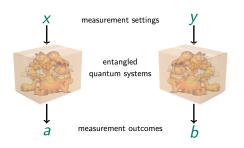
Quantum entanglement

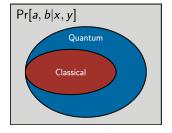
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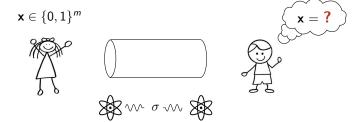


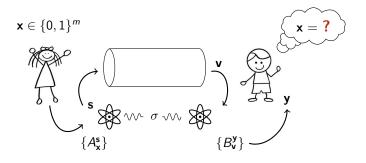
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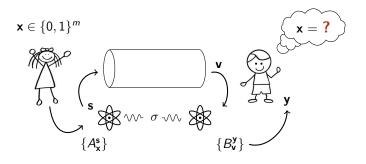








Can entanglement be beneficial?



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Depends

• Shannon regime: probability of error tends to zero as number of uses of the channel goes to infinity

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X

Finite number of channel uses

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• Finite number of channel uses

Prevedel et al.(2011): experimental result



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Finite number of channel uses

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Probability of error equal to zero

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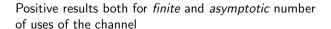


Finite number of channel uses

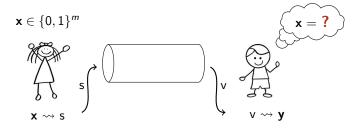
Prevedel et al.(2011): experimental result

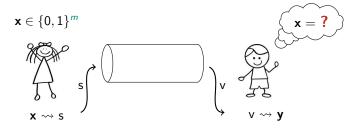


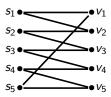
Probability of error equal to zero

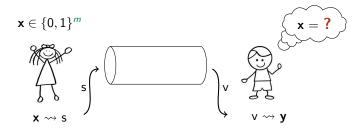


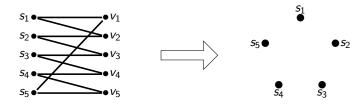


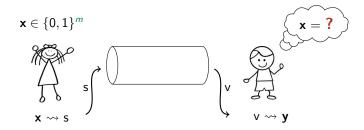


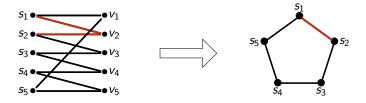


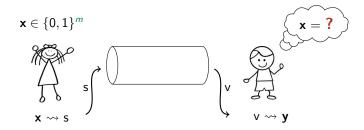


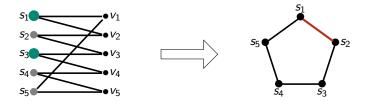


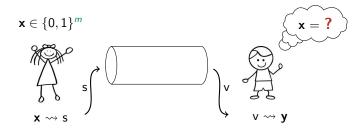


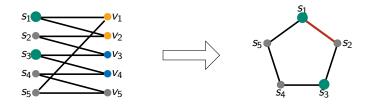












	One-shot	Asymptotic
Classical		
Quantum		

	One-shot	Asymptotic
Classical	$\alpha(G)$	<i>c</i> (<i>G</i>)
Quantum		

	One-shot		Asymptotic
Classical	$\alpha(G)$	<	<i>c</i> (<i>G</i>)
Quantum			

	One-shot		Asymptotic	
Classical	$\alpha(G)$	<	<i>c</i> (<i>G</i>)	Shannon 1956
Quantum				

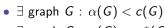
• \exists graph G : $\alpha(G) < c(G)$

	One-shot		Asymptotic
Classical	$\alpha(G)$	<u>≤</u>	c(G)
Quantum	$\alpha^{\star}(G)$		

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$$\exists$$
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	One-shot		Asymptotic
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	^		
Quantum	$\alpha^{\star}(G)$		

Cubitt et al. 2010



• \exists graph G : $\alpha(G) < \alpha^*(G)$



	One-shot		Asymptotic
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Quantum	$\alpha^{\star}(G)$	\leq	$c^{\star}(G)$

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	One-shot		Asymptotic
Classical	$\alpha(G)$	<u>≤</u>	c(G)
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Quantum	$\alpha^{\star}(G)$	\leq	$c^{\star}(G)$

- \exists graph G : $\alpha(G) < c(G)$
- \exists graph G : $\alpha(G) < \alpha^*(G)$

	One-shot		Asymptotic	
Classical	$\alpha(G)$	≤ 7	<i>c</i> (<i>G</i>)	Leung et al. 2012
Quantum	$\alpha^{\star}(G)$	<u>≤</u>	$c^*(G)$	Briët et al. 2012

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	One-shot		Asymptotic	
Classical	$\alpha(G)$	\leq	c(G)	
	\wedge I		\wedge	Briët et al. 2015
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Thank you!